## The geometric cobordism hypothesis Lecture 5a: The leftovers

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These slides: https://dmitripavlov.org/lecture-5a.pdf



- Tuesday: definitions
- Wednesday: locality and how to use it to prove one half of the GCH
- Thursday: the framed GCH (the other half)
- Today: contractibility of moduli spaces of cuts and its applications to GCH

#### Proposition (G.–P. (formal))

$$\begin{split} & \text{Given } d \geq 0, \text{ we have a left } Quillen \text{ functor} \\ & \mathrm{sPSh}(\mathsf{FEmb}_d)_{\mathsf{inj}} \to \mathrm{sPSh}(\mathsf{Cart} \times \Gamma \times \Delta^{\times d})_{\mathsf{loc}}, \qquad \mathcal{S} \mapsto \mathfrak{Bord}_d^{\mathcal{S}}. \end{split}$$

#### Theorem (G.–P.)

Given  $d \ge 0$ , the left derived functor of the left Quillen functor  $\mathrm{sPSh}(\mathsf{FEmb}_d)_{inj} \to \mathrm{sPSh}(\mathsf{Cart} \times \Gamma \times \Delta^{\times d})_{\mathsf{loc}}, \qquad S \mapsto \mathfrak{Bord}_d^S$ sends Čech nerves of open covers in  $\mathsf{FEmb}_d$  to weak equivalences.

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sends the Čech nerve of an open cover  $\{W_a \to U_a\}_{a \in A}$  of  $(W \to U) \in FEmb_d$  to a weak equivalence:

$$\operatorname{hocolim}_{n\in\Delta^{\operatorname{op}}}\coprod_{\alpha:[n]\to A}\mathfrak{Bord}_d^{W_\alpha\to U_\alpha}\xrightarrow{\sim}\mathfrak{Bord}_d^{W\to U},$$

where  $W_{\alpha} = W_{\alpha_0} \cap \cdots \cap W_{\alpha_n}$ .

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Step 1 Replace hocolim by colim Step 2 Pass to *n*-dimensional stalks on Cart for all  $n \ge 0$ . Step 3 Introduce a filtration (on *n*-dimensional stalks)

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Step 4 Prove all maps in the filtration are weak equivalences.

#### The codescent property: filtration

$$\operatorname{colim}_{n\in\Delta^{\operatorname{op}}}\coprod_{\alpha:[n]\to A}\mathfrak{Bord}_d^{W_\alpha\to U_\alpha}\to B_0\to\dots\to B_d\to\mathfrak{Bord}_d^{W\to U}.$$

#### Definition

Given 
$$d \ge 0$$
 and  $(W = \mathbf{R}^d \times U \to U) \in \mathsf{FEmb}_d^{\mathrm{op}}$ , the set  $\mathfrak{Bord}_d^{\mathbf{R}^d \times U \to U}(V, \langle \ell \rangle, \mathbf{m})_n$  has elements:

- a smooth manifold M;
- a V-family of embeddings  $M \to \mathbf{R}^d$ ;
- a  $V \times \Delta^n$ -family of cut tuples with  $m_1 \times \cdots \times m_d$  cells;
- $P: M \to \langle \ell \rangle;$
- smooth map  $V \to U$ ;

# $\operatorname{colim}_{n\in\Delta^{\operatorname{op}}}\coprod_{\alpha:[n]\to A}\mathfrak{Bord}_d^{W_\alpha\to U_\alpha}\to B_0\to\dots\to B_d\to\mathfrak{Bord}_d^{W\to U}.$

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#### Proposition (Formal)

The map  $\operatorname{colim} \to B_0$  is a weak equivalence in  $\operatorname{sPSh}(\Gamma \times \Delta^{\times d})_{\mathsf{loc}}$ .

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- B<sub>i</sub>: bordisms that can be chopped in the *i*th direction so that every piece belongs to B<sub>i-1</sub>.

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#### Proof.

• Evaluate  $B_{i-1} \to B_i$  on an arbitrary object X of  $\Gamma \times \Delta^{\times (d-1)}$ , obtaining a map  $B_{i-1}(X) \to B_i(X)$  in  $\mathrm{sPSh}(\Delta)$ ;

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- Extract the kth simplicial degree (for some k ≥ 0), obtaining a map in PSh(Δ) = sSet;
- The resulting simplicial set has
  - vertices: germs of cuts (embedded in W);
  - edges: bordisms between cuts (embedded in W);
  - 2-simplices: composition of bordisms;
  - everything is in smooth families indexed by  $\Delta^k$ ;
  - bordisms must belong to  $B_{i-1}$  respectively  $B_i$ .

Want to show:  $B_{i-1} \rightarrow B_i$  is a categorical weak equivalence in the Joyal model structure on simplicial sets.

## Intermission: Necklace categories

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- Answer: Dugger–Spivak necklace categories.

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# Intermission: Necklace categories of bordisms

•  $X = B_{i-1}$  or  $B_i$ , evaluated at  $X \in \Gamma \times \Delta^{\times (d-1)}$  and some  $[I] \in \Delta$  (smooth families of bordisms indexed by  $\Delta^I$ ).

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- Observation: the ambient composed bordism never changes ⇒ can fix it in advance.

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- This implies  $B_{i-1} \rightarrow B_i$  is a weak equivalence.

# The big picture

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- All neighborhoods can be chosen to be subordinate to the open cover of W.
- How does this help us to show contractibility of necklace categories?

A Kan complex X is contractible if and only if any map  $\partial \Delta^n \to X$  can be simplicially homotoped to a map that extends along  $\partial \Delta^n \to \Delta^n$ .

• Apply this theorem to the fibrant replacement X of the nerve of the necklace category of  $B_{i-1}$  (or  $B_i$ ) from x to y.

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- Chop up *M* as explained on the previous slide.

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- Chop up *M* as explained on the previous slide.
- By induction on the Morse decomposition, push the cuts past each small region in the Morse decomposition, with some cutting and gluing of cuts.
- At the final step, all cuts have been collapsed to the source cut of *M*.

# How is this used in the framed GCH?

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- base:  $\mathbf{R}$  Map( $\mathfrak{Bord}_d^{\mathbf{R}^d \times U \to U}, \mathcal{V}$ ) =  $\mathbf{R}$  Map( $B_d, \mathcal{V}$ );
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- everything except for the index k handle is in  $B_{k-1}$ ;
- the value on the index k handle is unique up to a contractible choice;
- hence,  $\mathbf{R} \operatorname{Map}(B_k, \mathcal{V}) \simeq \mathbf{R} \operatorname{Map}(B_{k-1}, \mathcal{V})_{\operatorname{unit}}$ ;

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# Cutting out handles

#### Theorem (G.–P.)

For any  $d \ge 0$  and  $0 \le k \le d$ , the following squares are homotopy cocartesian in  $C^{\infty}Cat_{\infty,d}^{\otimes,\vee}$ :  $O_{k-1} \longrightarrow \overline{O}_{k-1} \longrightarrow \widetilde{O}_{k-1} \longrightarrow B_{k-1}$  $\dot{H}_k \longrightarrow \overline{H}_k \longrightarrow \widetilde{H}_k \longrightarrow \widetilde{B}_k.$ ■ *H<sub>k</sub>*: index *k* − 1 (counit) and index *k* (unit) handles; •  $O_{k-1}$ : index k-1 (counit) handle; **•**  $\overline{H}$ ,  $\overline{O}$ : same, with tails attached in the (d-2)nd direction; • H, O: same, with tails attached in the (d-1)st direction. left two squares: insert cuts close to the handle;

 right square: invoke the same proof as for locality, using a new open cover.

# Nebulous visions of the future...

- Prequantum/classical: the book "Natural operations in differential geometry" (Kolář–Michor–Slovák) constructs a lot of functorial field theories...
- Quantum: quantization and path integrals (for fully extended FFTs) via the GCH;
- Further compute the right hand side of GCH via  $\infty$ -Lie theory;

Explore possible value categories:

- geometric factorization algebras (Peña, based on Carmona–Flores–Muro);
- closed symmetric monoidal category with duals of complete vector spaces (?).